

# A FTPL Approach to Online Non-convex Optimization

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CSIE 5002 – Prediction, Learning, and Games

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# Convexity

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# Why is convexity important?

## Convexity:

- Ensures that **the local minimum is the global minimum.**
- Optimality condition helps capture the minimizer.

## Definition (Non-convex)

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We say a problem is **convex** if **both** the objective function  $f$  and the constraint set  $\mathcal{X}$  are convex; otherwise, we say it is a **non-convex** problem.

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# From Convex to Non-convex

## Convex:

- Convexity ensures **the global minimum**.

## Non-convex:

- There exists potentially many local minima.
- Finding the global minima among all the local minima is hard (at least NP-Hard). <sup>1</sup>
- Saddle points present.

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<sup>1</sup>See an overview by [Danilova et al. \(2021\)](#).

# Challenge in Online Non-convex Optimization

- Suggala and Netrapalli (2020) prove that no deterministic algorithm can achieve  $o(1)$  regret in the setting of online non-convex learning.
- The broadly developed methods such as Follow the Regularized Leader (FTRL) and Online Mirror Descent (OMD) cannot perform well in the non-convex setting.

# A Simple Neural Network (NN) Example

## Algorithm (Model Training in Classification)

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Given training data  $(x, y) \in (\mathcal{X}, \mathcal{Y})$ , the activation function, and the loss function, LEARNER announces a weight  $w_t \in \mathcal{W}$  corresponding to model structure at  $t$ -th epoch and receives a cost  $f: (\mathcal{X}, \mathcal{Y}) \times \mathcal{W} \rightarrow \mathbb{R}$ .

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- Imagine an example where the activation function is **Sigmoid**, which is non-convex to  $w \in \mathcal{W}$ .

# Non-convex Applications

## Neural Networks:

- Szegedy et al. (2013) specify loss functions for neural networks are non-convex in general.
- Applications include adversarial training (Generative Adversarial Networks, GAN), efficient computation of an equilibrium in non-convex games, etc.

# Optimization

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# Reduce Regret Minimization to Optimization

## Online-to-offline reduction:

- [Kalai and Vempala \(2005\)](#) introduce Follow the Perturbed Leader (FTPL) to the regret minimization problem.
- The action set and loss functions for LEARNER is **not** assumed to be convex.
- The minimizer is assumed to be computed **efficiently**.

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Some heuristic methods can find approximate global optima reasonably fast. See [Drori and Shamir \(2020\)](#).

# Optimization Oracle

## Definition (Value Oracle)

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A **value oracle** is a procedure  $\text{Val} : \mathcal{X} \times \mathcal{Y} \rightarrow [0, 1]$  that for any action pair  $x \in \mathcal{X}, y \in \mathcal{Y}$ , return the loss value  $\ell(x, y)$  in time  $\mathcal{O}(1)$ .

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## Definition (Optimization Oracle)

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An **optimization oracle** is a procedure  $\text{Opt}$  that receives a distribution  $q \in \Delta(\mathcal{Y})$  and returns the best performing action w.r.t.  $q$  in  $\mathcal{O}(1)$  on any input, i.e.,

$$\forall q \in \Delta(\mathcal{Y}), \text{Opt}(q) \in \arg \min_{x \in \mathcal{X}} \mathbb{E}_{y \sim q}[\ell(x, y)].$$

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# Efficient Computation Assumption

## Online with offline:

- [Littlestone \(1989\)](#) shows that efficient online computation implies efficient offline approximation.
- What [Kalai and Vempala \(2005\)](#) and we discuss here is how to use efficient offline optimization algorithms for the online problem.

**FTPL**

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# Preliminaries

## Setting and assumptions:

- $\mathcal{X} \subseteq \mathbb{R}^d$  is the action set for LEARNER which is bounded with  $\ell_\infty$  diameter of  $D$ , i.e.,  $D = \sup_{x,y \in \mathcal{X}} \|x - y\|_\infty$ .
- The sequence of loss function  $f_t$  chosen by REALITY are  $L$ -Lipschitz w.r.t.  $\ell_1$  norm, i.e.,  $|f_t(x) - f_t(y)| \leq L\|x - y\|_1$  for all  $x, y \in \mathcal{X}$ .
- LEARNER aims to choose a sequence of actions  $\{x_t\}_{t=1}^T$  to minimize the regret

$$R_T(x) = \frac{1}{T} \sum_{t=1}^T f_t(x_t) - \frac{1}{T} \inf_{x \in \mathcal{X}} \sum_{t=1}^T f_t(x).$$

# Recall of Follow the Leader (FTL)

## Algorithm (Follow the Leader, FTL)

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LEARNER plays a strategy  $x_t \in \mathcal{X}$  at  $t$ -th round with minimal loss over past rounds, i.e.,

$$x_t = \arg \min_{x \in \mathcal{X}} \sum_{i=1}^{t-1} f_i(x).$$

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# Follow the **Perturbed** Leader (FTPL)

## Randomness:

- Following FTL, we perturb the cumulative loss by **adding a random perturbation** at each round.

## Algorithm 1 (FTPL, Agarwal et al. (2019))

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Given a parameter  $\eta > 0$ , LEARNER draws an i.i.d. random vector  $\sigma_t \sim (\exp(\eta))^d$  and plays a strategy

$$x_t = \arg \min_{x \in \mathcal{X}} \left\{ \sum_{i=1}^{t-1} f_i(x) - \sigma_t^\top x \right\}$$

at  $t$ -th round.

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# Analysis

## Complexity:

- **Algorithm 1** achieves the regret bound  $\mathbb{E}[R_T(x)] \leq \mathcal{O}(T^{-1/3})$ .
- **Cesa-Bianchi and Lugosi (2006)** show that any algorithm that is guaranteed to work against an oblivious REALITY also works for a non-oblivious REALITY. That is, it suffices to work with a single random vector  $\sigma$ .



# Strategy as an Oracle

LEARNER'S strategy can be regarded as an offline optimization oracle.

## Definition (Offline optimization oracle)

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An offline optimization oracle is a procedure `Off-Opt` that receives a sequence of loss functions  $f_1, \dots, f_k$  and a  $d$ -dimensional vector  $\sigma$  and returns the best performing action

$$x^* = \arg \min_{x \in \mathcal{X}} \left\{ \sum_{i=1}^k f_i(x) - \sigma^\top x \right\}.$$

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# Approximate Optimization Oracle

## Definition $((\alpha, \beta)$ -approximate optimization oracle)

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An optimization oracle is called  $(\alpha, \beta)$ -approximate optimization oracle (**Approx-Opt**) if it returns  $x^* \in \mathcal{X}$  such that

$$f(x^*) - \sigma^\top x^* \leq \inf_{x \in \mathcal{X}} \{f(x) - \sigma^\top x\} + (\alpha + \beta \|\sigma\|_1).$$

We denote it with **Approx-Opt** $_{\alpha, \beta}(f - \langle \sigma, \cdot \rangle)$ .

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# FTPL with Approximate Optimization Oracle

## Algorithm 2 (FTPL, Suggala and Netrapalli (2020))

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Given a parameter  $\eta > 0$ , LEARNER draws an i.i.d. random vector  $\sigma_t \sim (\exp(\eta))^d$  and plays a strategy

$$x_t = \text{Approx-Opt}_{\alpha, \beta} \left( \sum_{i=1}^{t-1} f_i - \langle \sigma_t, \cdot \rangle \right)$$

at  $t$ -th round.

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# Performance Analysis

- Suggala and Netrapalli (2020) give the regret bound

$$\mathbb{E}[R_T(x)] \leq \mathcal{O}\left(\eta d^2 D L^2 + \frac{d(\beta T + D)}{\eta T} + \alpha + \beta d L\right).$$

- By setting appropriate  $\eta = \left(\sqrt{dT} - L\right)^{-1}$  and when  $\alpha = \mathcal{O}(T^{-1/2})$ ,  $\beta = \mathcal{O}(T^{-1})$ , FTPL achieves the  $\mathcal{O}(T^{-1/2})$  regret bound.

**OFTPL**

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# Online Learning with Predictable Sequences

## Optimistic case:

- The general online learning setting assumes that the loss functions are chosen in an adversarial manner by REALITY.
- However, the loss functions might have some patterns and could be predictable.
- LEARNER can have **prior knowledge** about the loss functions (see **Rakhlin and Sridharan (2013)**).

# OFTPL with Approximate Optimization Oracle

## Algorithm 3 (OFTPL, Suggala and Netrapalli (2020))

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Let  $g_t$  be LEARNER'S guess of  $f_t$  depending on  $f_1, \dots, f_{t-1}$ . Given a parameter  $\eta > 0$ , LEARNER draws an i.i.d. random vector  $\sigma_t \sim (\exp(\eta))^d$  and plays a strategy

$$x_t = \text{Approx-Opt}_{\alpha, \beta} \left( \sum_{i=1}^{t-1} f_i + g_t - \langle \sigma_t, \cdot \rangle \right)$$

at  $t$ -th round.

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Rakhlin and Sridharan (2013) give a thorough discussion on the choices of  $g_t$ .

# Performance Analysis

- Suggala and Netrapalli (2020) give the regret bound

$$\mathbb{E}[R_T(x)] \leq \mathcal{O}\left(\eta d^2 D \sum_{t=1}^T \frac{L_t^2}{T} + \frac{d(\beta T + D)}{\eta T} + \alpha + \beta d \sum_{t=1}^T \frac{L_t}{T}\right),$$

where they assume  $g_t - f_t$  is  $L_t$ -Lipschitz w.r.t.  $\ell_1$  norm.

- The regret bound is expected to be smaller when the guess  $g_t$  is close to  $f_t$ .



# Summary

## Takeaway:

- Non-convex problems are common and hard to solve, but a slightly modified FTL can achieve a sublinear regret bound.
- The Exponential Weight Algorithm is another approach to solve non-convex problems.

## Comment:

- Can other distributions improve the regret bound?

# Reference

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